

A Lower Bound for Potentially F -Graphic Degree Sequences

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The degree sequence problem

Problem: Given an integer sequence $\mathbf{d} = (d_1, \dots, d_n)$ determine if there exists a graph G with \mathbf{d} as its sequence of degrees.

If such a G exists then \mathbf{d} is said to be *graphic*, and G is called a *realization*.

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Havel (1955) and Hakimi (1962) gave an algorithm to decide.

$$(3, 3, 3, 3, 3, 3) \rightarrow (2, 2, 2, 3, 3) = (3, 3, 2, 2, 2) \rightarrow (2, 1, 1, 2) = (2, 2, 1, 1) \rightarrow (1, 0, 1) = (1, 1, 0) \rightarrow (0, 0)$$

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To construct a realization, work backwards using simple edge augmentations.

Erdős-Gallai criterion

Theorem

[Erdős, Gallai (1960)]

A nonincreasing sequence of nonnegative integers $\mathbf{d} = (d_1, \dots, d_n)$ ($n \geq 2$) is graphic if, and only if, $\sum_{i=1}^n d_i$ is even and for each integer k , $1 \leq k \leq n - 1$,

$$\sum_{i=1}^k d_i \leq k(k-1) + \sum_{i=k+1}^n \min\{k, d_i\}.$$

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The degrees of the first k vertices are “absorbed” within k -subset and the degrees of remaining vertices. A necessary condition which is also sufficient!

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For a graphic \mathbf{d} , $\sum_{i=1}^n d_i \geq 2(n-1)$ if and only if there exists a connected G realizing \mathbf{d} .

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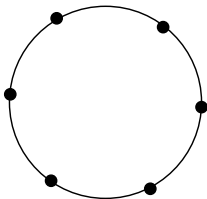
Proof: (Sufficiency) If there exists a connected realization then G contains a spanning tree. Thus G has $n-1$ edges and so $\sum_{i=1}^n d_i \geq 2(n-1)$.

(Necessity) Pick the realization of \mathbf{d} with the fewest number of components. If this number is 1, then we are done. Otherwise one of the components contains a cycle. Performing a simple edge-exchange allows us to move to a realization with fewer components. \square

For a subgraph F , \mathbf{d} is said to be **potentially F -graphic** if there exists a realization of \mathbf{d} containing F .

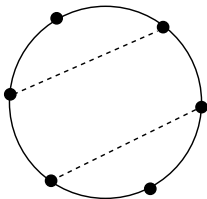
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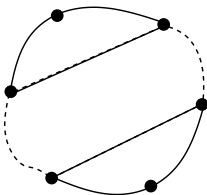
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Problem

Given a subgraph F , determine the *least* even integer m s.t.
 $\sum d_i \geq m \Rightarrow \mathbf{d}$ is potentially F -graphic.

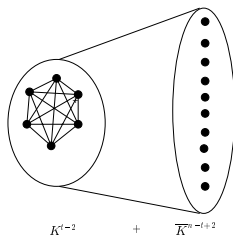
Denote m by $\sigma(F, n)$.

Erdős, Jacobson, Lehel Conjecture

Conjecture

(EJL - 1991) For n sufficiently large,
 $\sigma(K^t, n) = (t - 2)(2n - t + 1) + 2$.

Lower bound arises from considering:



$$\mathbf{d} = ((n - 1)^{t-2}, (t - 2)^{n-t+2})$$

Erdős, Jacobson, Lehel Conjecture

Conjecture settled:

- ▶ $t = 3$ Erdős, Jacobson, & Lehel(1991),
- ▶ $t = 4$ Gould, Jacobson, & Lehel(1999), Li & Song(1998),
- ▶ $t = 5$ Li & Song(1998),
- ▶ $t \geq 6$ Li, Song, & Luo(1998)
- ▶ $t \geq 3$ S.(2005), Ferrara, Gould, S. (2009) - purely graph-theoretic proof.

Theorem

For n sufficiently large, $\sigma(K^t, n) = (t - 2)(2n - t + 1) + 2$.

Sketch of our proof

- ▶ Uses induction on t .
- ▶ Erdős-Gallai guarantees enough vertices of high degree.
- ▶ Uses notion of an edge-exchange.
- ▶ Edge-exchange allows us to place desired subgraph on vertices of highest degree and “build” K^t from smaller clique guaranteed by inductive hypothesis.

Extending the EJL-conjecture to an arbitrary graph F

Let F be a forbidden subgraph.

Let $\alpha(F)$ denote the independence number of F and define:

$$u := u(F) = |V(F)| - \alpha(F) - 1,$$

and

$$s := s(F) = \min\{\Delta(H) : H \subset F, |V(H)| = \alpha(F) + 1\}.$$

Consider the following sequence,

$$\pi(F, n) = ((n-1)^u, (u+s-1)^{n-u}).$$

An example

Consider $F = K_{6,6}$. Then,

$$u(K_{6,6}) = |V(K_{6,6})| - \alpha(K_{6,6}) - 1 = 12 - 6 - 1 = 5$$

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Thus,

$$\pi(K_{6,6}, n) = ((n-1)^5, (5+4-1)^{n-5}) = ((n-1)^5, 8^{n-5}).$$

A General Lower Bound

If F' is a subgraph of F then $\sigma(F', n) \leq \sigma(F, n)$ for every n . Let $\sigma(\pi)$ denote the sum of the terms of π .

Proposition (Ferrara, S. - '09)

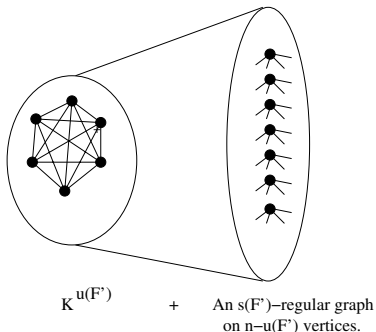
Given a graph F and n sufficiently large then,

$$\sigma(F, n) \geq \max\{\sigma(\pi(F', n)) + 2|F' \subseteq F\} \quad (1)$$

$$= \max\{n(2u(F') + s(F') - 1) | F' \subseteq F\} \quad (2)$$

Proof of Lower Bound

PROOF: Let $F' \subseteq F$ be the subgraph which achieves the max.
 Consider,



$$u(F') = |V(F')| - \alpha(F') - 1$$

$$s(F') = \min\{\Delta(H) : H \subset F', |V(H)| = \alpha(F') + 1\}$$

Let's do a little better with our Example

Let $F = K_{6,6}$.

Then $u(K_{6,6}) = 12 - 6 - 1 = 5$ and $s(K_{6,6}) = 4$.

Consider,

$$\pi^*(K_{6,6}, n) = ((n-1)^5, 10, 9, 8^{n-7}).$$

Allowing a few vertices to have a little higher degree

Let $v_i(H)$ be the number of vertices of degree i in H . Let $M_i(H)$ denote the set of induced subgraphs on $\alpha + 1$ vertices with $v_i(H) > 0$.

For all i , $s \leq i \leq \alpha - 1$ define:

$$m_i = \min_{M_i(H)} \{\text{vertices of degree at least } i\}$$

$$n_s = m_s - 1 \text{ and } n_i = \min\{m_i - 1, n_{i-1}\}$$

Finally, set $\delta_{\alpha-1} = n_{\alpha-1}$ and for all i , $s \leq i \leq \alpha - 2$ define $\delta_i = n_i - n_{i+1}$ and

$$\begin{aligned} \pi^*(F, n) = & ((n-1)^u, (u+\alpha-1)^{\delta_{\alpha-1}}, (u+\alpha-2)^{\delta_{\alpha-2}}, \dots \\ & (u+s)^{\delta_s}, (u+s-1)^{n-u-\sum \delta_i}). \end{aligned}$$

An Example

Let $F = K_{6,6}$.

Then $u(K_{6,6}) = 12 - 6 - 1 = 5$ and $s(K_{6,6}) = 4$.

$m_4 = 3$ and $m_5 = 2$

$n_4 = m_4 - 1 = 2$ and $n_5 = \min\{m_5 - 1, n_4\} = 1$

$\delta_5 = n_5 = 1$ and $\delta_4 = n_5 - n_4 = 1$

Thus,

$$\pi^*(K_{6,6}, n) = ((n-1)^5, 10, 9, 8^{n-7})$$

A Stronger Lower Bound

Theorem (Ferrara, S. - '09)

Given a graph F and n sufficiently large then,

$$\sigma(F, n) \geq \max\{\sigma(\pi^*(F', n)) + 2|F' \subseteq F\}$$

When Does Equality Hold?

- ▶ **cliques**
- ▶ **complete bipartite graphs** Chen, Li, Yin '04; Gould, Jacobson, Lehel '99; Li, Yin '02
- ▶ **complete multipartite graphs** Chen, Yin '08; G. Chen, Ferrara, Gould, S. '08; Ferrara, Gould, S. '08
- ▶ **matchings** Gould, Jacobson, Lehel '99
- ▶ **cycles** Lai '04
- ▶ **(generalized) friendship graph** Ferrara, Gould, S. '06, (Chen, S., Yin '08)
- ▶ **clique minus an edge** Lai '01; Li, Mao, Yin '05
- ▶ **disjoint union of cliques** Ferrara '08

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(weaker version) Given a graph F , let $\epsilon > 0$. Then there exists an $n_0 = n_0(\epsilon, F)$ such that for any $n > n_0$

$$\sigma(F, n) \leq \max\{(n(2u(F')) + d(F') - 1 + \epsilon) \mid F' \subseteq F\}.$$

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Conjecture (strong form) holds for **graphs with independence number 2** (Ferrara, S. - '09)

An example of the generalized problem

Is the following graphic?

$$\langle \mathbb{V}, \mathbf{d}, D \rangle = \langle \{V_1, V_2\}, (5^4, 3^8), \begin{bmatrix} 6 & 8 \\ 8 & 8 \end{bmatrix} \rangle$$

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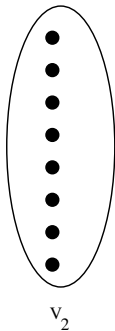
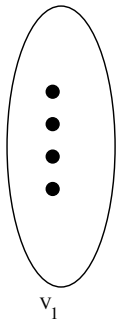
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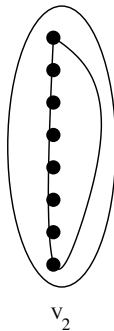
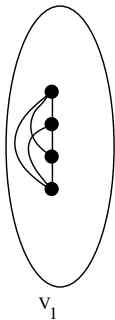
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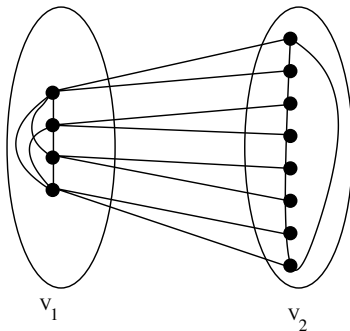
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Let $\mathbf{d} = (d_1^{v_1}, d_2^{v_2}, \dots, d_k^{v_k})$ where $v_i = |V_i|$ and so V_i is the set of vertices of degree d_i . Let $\mathbb{V} = \{V_1, \dots, V_k\}$. Let $D = (d_{ij})$ be a $k \times k$ matrix, with d_{ij} denoting the number of edges between V_i and V_j ; d_{ii} is the number of edges contained entirely within V_i .

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Joint degree-matrix graphic realization problem

Given $\langle \mathbb{V}, \mathbf{d}, D \rangle$, decide whether a simple graph G exists such that, for all i , each vertex in V_i has degree d_i , and, for $i \neq j$, there are exactly d_{ij} edges between V_i and V_j , while, for all i , there are exactly d_{ii} edges contained in V_i .

Amanatidis, Green and Mihail (AGM) have shown that the following natural necessary conditions for a realization to exist are also sufficient. The conditions are:

Degree feasibility: $2d_{ii} + \sum_{j \in [k], j \neq i} d_{ij} = v_i d_i$, for all $1 \leq i \leq k$, and

Matrix feasibility: D is a symmetric matrix with non-negative integral entries, $d_{ij} \leq v_i v_j$, for all $1 \leq i \leq k$, and $d_{ij} \leq \binom{v_i}{2}$, for all $1 \leq i \leq k$.

Theorem (Joint Degree-Matrix Realization Theorem - AGM)

Given $\langle \mathbb{V}, \mathbf{d}, D \rangle$, if degree and matrix feasibility hold, then a graph G exists that realizes $\langle \mathbb{V}, \mathbf{d}, D \rangle$. Furthermore, such a graph can be constructed in time polynomial in n .

Can one prove our conjectures using the Joint Degree-Matrix Theorem?

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The value of d_{11} forces, by degree feasibility, $d_{12} = d_{21} = 30$. In turn, by degree feasibility again, we get $d_{22} = 1$. It is now easy to check that matrix feasibility holds. Thus, \mathbf{d} has a realization containing a copy of K^6 .

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- ▶ Proved the EJM-conjecture
- ▶ Generalized the EJM-conjecture and proved a specific case
- ▶ Joint Degree-Matrix Theorem appears to be a useful tool.
- ▶ Can we use it to prove the generalized EJM-conjecture?

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Thank you!